MUSE: Mapping Understanding and deSign by Example

Bogdan AlexeUC Santa CruzLaura ChiticariuUC Santa CruzRenée J. MillerU. of TorontoWang-Chiew TanUC Santa Cruz

April 8, 2008







One of the first steps in information integration is to specify the relationships (schema mappings) between schemas. This is known to be a difficult task.



Example: Data Exchange



- I Mapping systems can automate only part of the mapping process
 - Typically, intricate manual work is needed to perfect the specification

- Mapping systems can automate only part of the mapping process
 - Typically, intricate manual work is needed to perfect the specification
 - The visual specification may be ambiguous. Mapping systems make default choices to resolve the ambiguities
 - These choices may not correspond to a designer's intentions
 - The mapping designer might refine the specification manually

Real Life Example

In real life scenarios, mappings are extremely complicated



Real Life Example

In real life scenarios, mappings are extremely complicated

Map 2:			
	for sm	2x0 in S0.dummy COUNTRY 4	
	exists	tm2x0 in S27.dummy country 10, tm2x1 in S27.dummy organiza 13	
	when	<pre>tm2x0.country.membership=tm2x1.organization.id,</pre>	
	satisf	<pre>sm2x0.COUNTRY.AREA=tm2x0.country.area, sm2x0.COUNTRY.CAPITAL=tm2x0.country.capital,</pre>	
		<pre>sm2x0.COUNTRY.CODE=tm2x0.country.id, sm2x0.COUNTRY.NAME=tm2x0.country.name,</pre>	
		<pre>sm2x0.COUNTRY.POPULATION=tm2x0.country.population,(</pre>	
	Map 3:		
	100000000000000000000000000000000000000	for sm3x0 in S0.dummy_GE0_RIVE_23, sm3x1 in S0.dummy_RIVER_24,	
		where sm3x0.GE0_RIVER.RIVER=sm3x1.RIVER.NAME, sm3x2.PROVINCE.NAME=sm3x0.GE0_RIVER.PROVINCE,	
		smiskz.rnovince.countri-smizko.countri.cobe,	
		track in S27. dummy_live44, that i that i that dummy_live_cated_23,	
		tm3x4 in 527. dummy_organiza 13	
		tm3x0 in 327.dummy_organiza_13	
		tm2v6 country.id=tm3v1 located country	
		satisf sm2x0. COUNTRY AREA=tm3x4 country area sm2x0. COUNTRY CAPITAL =tm3x4 country capital	
		Sm2x0 COUNTRY CODE=tm3x4 country id sm2x0 COUNTRY NAME=tm3x4 country name	
		sm2x0.COUNTRY.POPULATION=tm3x4.country.population.sm3x1.BIVER.LENGTH=tm3x0.river.length.	
		sm3x0.GE0 RIVER.COUNTRY=tm3x1.located.country.sm3x0.GE0 RIVER.PROVINCE=tm3x1.located.province.	
		sm3x1.RIVER.NAME=tm3x0.river.name).(
	Map 4:		
	1	for sm4x0 in S0.dummy GEO ISLA 25, sm4x1 in S0.dummy ISLAND 26,	
		sm4x2 in S0.dummy PROVINCE 5	
		where sm4x0.GE0 ISLAND.ISLAND=sm4x1.ISLAND.NAME, sm4x2.PROVINCE.NAME=sm4x0.GE0 ISLAND.PROVINCE,	
		<pre>sm4x2.PROVINCE.COUNTRY=sm2x0.COUNTRY.CODE,</pre>	
		exists tm4x0 in S27.dummy_island_26, tm4x1 in tm4x0.island.dummy_located_25,	
		tm4x4 in S27.dummy_country_10, tm4x5 in tm4x4.country.dummy_province_9,	
		tm4x6 in S27.dummy_organiza_13	
		where tm4x4.country.membership=tm4x6.organization.id, tm4x5.province.id=tm4x1.located.province,	
		tm2x0.country.id=tm4x1.located.country,	
		<pre>satisf sm2x0.COUNTRY.AREA=tm4x4.country.area, sm2x0.COUNTRY.CAPITAL=tm4x4.country.capital,</pre>	
		sm2x0.COUNTRY.CODE=tm4x4.country.id, sm2x0.COUNTRY.NAME=tm4x4.country.name,	
		sm2x0.COUNTRY.POPULATION=tm4x4.country.population, sm4x1.ISLAND.AREA=tm4x0.island.area,	
		<pre>sm4x1.ISLAND.COORDINA1ESLA1=tm4x0.island.latitude, sm4x0.GE0_ISLAND.COUNTRY=tm4x1.located.country,</pre>	
		sm4x0.GEO_ISLAND.PROVINCE=Tm4x1.located.province, sm4x1.ISLAND.COORDINATESLONG=tm4x0.island.longitud	le,
	И	SM4XI.ISLAND.NAME=TM4XU.ISland.name),(
	Map 5:		
		TOT Smoxu in Suldummy_GEU_SEA_19, Smox1 in Suldummy_SEA_20,	
		SMDXZ IN SUJUMUY PROVINCE O SEA DROVINCE OFFYA CEO SEA SEA-OFFYI SEA NAME	
		SHISKZ-FROWINCE-COUNTRI-SHIZAO-COUNTRI-CODE	
		tm5x4 in 527. dummy_sea_19, tm5x1 in tm5x4.comptor dummy_service 9	
		tm5x4 in 527. dummy_organiza 13	
		where they's country membership=th5% organization id th5%5 province id=th5%1 located province	
		tm2x6 country id=tm5x1 located country	
		satisf sm2x0 COUNTRY AREA=tm5x4 country area sm2x0 COUNTRY CAPITAL=tm5x4 country capital	
		sm2x0.COUNTRY.CODE=tm5x4.country.id. sm2x0.COUNTRY.NAME=tm5x4.country.name	
		sm2x0.COUNTRY.POPULATION=tm5x4.country.population.sm5x1.SEA.DEPTH=tm5x0.sea.depth.	

sm5x0.GE0_SEA.COUNTRY=tm5x1.located.country, sm5x0.GE0_SEA.PROVINCE=tm5x1.located.province,

sm5x1.SEA.NAME=tm5x0.sea.name),(

- Specifications are often impossible to understand through visual inspection
- Few tools are available to assist in understanding and designing alternative mappings

- Specifications are often impossible to understand through visual inspection
- Few tools are available to assist in understanding and designing alternative mappings
 - MUSE is a tool designed towards this end

- Specifications are often impossible to understand through visual inspection
- Few tools are available to assist in understanding and designing alternative mappings
- MUSE is a tool designed towards this end
- In MUSE, we focus on declarative specifications



Our vision

- MUSE is a mapping design wizard that uses (real) data examples to help designers understand, design and refine schema mappings
- MUSE leverages familiar data examples to help understand mappings
 - real data examples are used whenever possible
 - otherwise, synthetic examples are constructed

Our vision

- MUSE is a mapping design wizard that uses (real) data examples to help designers understand, design and refine schema mappings
- MUSE leverages familiar data examples to help understand mappings
 - real data examples are used whenever possible
 - otherwise, synthetic examples are constructed
 - Currently, MUSE has two features
 - Muse-G: design grouping semantics
 - Muse-D: disambiguate alternative mappings

MUSE Workflow





























CompDB: Rcd **Companies:** Set of **Company:** *Rcd* ➤ cbranch cname location **Projects:** Set of **Project:** *Rcd* pid pname cbranch manager **Employees:** Set of **Employee:** *Rcd* → eid ename contact

CompDB: Rcd **Companies:** Set of **Company:** *Rcd* \rightarrow cbranch cname location **Projects:** Set of **Project:** *Rcd* pid pname cbranch manager **Employees:** Set of **Employee:** *Rcd* 🔶 eid ename contact

OrgDB: *Rcd* **Orgs:** Set of **Org:** *Rcd* oname Projects: Set of **Project:** *Rcd* pname manager -**Employees:** Set of **Employee:** *Rcd* eid \leftarrow ename

CompDB: Rcd **OrgDB:** *Rcd* **Companies:** Set of **Orgs:** Set of **Company:** *Rcd* **Org:** *Rcd* ➤ cbranch oname Projects: Set of cname location **Project:** *Rcd* pname **Projects:** Set of manager -**Project:** *Rcd* Employees: Set of pid **Employee:** *Rcd* pname eid < cbranch ename manager **Employees:** Set of **Employee:** *Rcd* → eid ename contact



Declarative Mapping

for

c in CompDB.Companies p in CompDB.Projects

e in CompDB.Employees

CompDB: Rcd **OrgDB:** *Rcd* **Orgs:** Set of **Companies:** Set of **Company:** *Rcd* **Org:** *Rcd* cbranch oname **Projects:** Set of cname location **Project:** *Rcd* pname **Projects:** Set of manager -**Project:** *Rcd* **Employees:** Set of pid **Employee:** *Rcd* pname eid < cbranch ename manager **Employees:** Set of **Employee:** *Rcd* eid ename contact

Declarative Mapping

for

- c in CompDB.Companies
- p in CompDB.Projects
- e in CompDB.Employees

satisfy

p.cbranch = c.cbranch e.eid = p.manager


Declarative Mapping

c in CompDB.Companies p in CompDB.Projects e in CompDB.Employees

p.cbranch = c.cbranche.eid = p.manager

- o in OrgDB.Orgs
- p₁ in o.Projects
- e₁ in OrgDB.Employees

CompDB: Rcd **OrgDB:** *Rcd* for **Orgs:** Set of **Companies:** Set of **Company:** *Rcd* **Org:** *Rcd* ➤ cbranch oname **Projects:** Set of cname location **Project:** *Rcd* pname **Projects:** Set of manager -**Project:** *Rcd* **Employees:** Set of pid **Employee:** *Rcd* pname eid cbranch ename manager **Employees:** Set of **Employee:** *Rcd* ≻ eid ename contact

Declarative Mapping

c in CompDB.Companies p in CompDB.Projects e in CompDB.Employees

satisfy

p.cbranch = c.cbranche.eid = p.manager

exists

o in OrgDB.Orgs

 p_1 in o.Projects

e₁ in OrgDB.Employees

satisfy

 p_1 manager = e_1 eid

CompDB: Rcd **OrgDB:** *Rcd* **Orgs:** Set of **Companies:** Set of **Company:** *Rcd* **Org:** *Rcd* ➤ cbranch oname **Projects:** Set of cname location **Project:** *Rcd* pname **Projects:** Set of manager _ **Project:** *Rcd* **Employees:** Set of pid **Employee:** *Rcd* pname eid < cbranch ename manager **Employees:** Set of **Employee:** *Rcd* ≻ eid ename contact

Declarative Mapping

for

c in CompDB.Companies p in CompDB.Projects e in CompDB.Employees **satisfy** p.cbranch = c.cbranch e.eid = p.manager

exists

- o in OrgDB.Orgs p₁ in o.Projects e₁ in OrgDB.Employees satisfy
 - p_1 manager = e_1 eid

where

c.cname = o.oname e.eid = e_1 .eid e.ename = e_1 .ename p.pname = p_1 .pname









Group by what subset of {cbranch, cname, location} ?

■ **Goal:** infer a grouping function that has the same effect as the one intended by the designer

- **Goal:** infer a grouping function that has the same effect as the one intended by the designer
 - Muse-G probes each possible grouping attribute: start with cbranch

- **Goal:** infer a grouping function that has the same effect as the one intended by the designer
- Muse-G probes each possible grouping attribute: start with cbranch

Example source

Companies

Redmond	Microsoft	USA
S. Valley	Microsoft	USA

Projects

Ρ1	DB	Redmond	e4
P2	Web	S. Valley	e5

Employees

e4	John	x234
e5	Anna	x888

- **Goal:** infer a grouping function that has the same effect as the one intended by the designer
 - Muse-G probes each possible grouping attribute: start with cbranch

Example source		Target Scenario 1 group by cbranch	Target Scenario 2 do not group by cbranch
Companies		Orgs	Orgs
Redmond Micros	oft USA	Microsoft	Microsoft
S. Valley Micros	oft USA	Projects:	Projects:
Projects		DB e4	DB e4
P1 DB Redmo	ond e4	Microsoft	Web e5
P2 Web S. Val	ley e5	Projects:	
Employees	5	Web e5	
e4 John x234		Employees	Employees
e5 Anna x888		e4 John	e4 John
		e5 Anna	e5 Anna

- **Goal:** infer a grouping function that has the same effect as the one intended by the designer
 - Muse-G probes each possible grouping attribute: start with cbranch

	Exa	mple	source		Target Scenario 1 group by cbranch	Target Scenario 2 do not group by cbranch
Cor	npan	ies			Orgs	Orgs
	Red	mond	Microsoft	USA	Microsoft	Microsoft
	S. ∖	/alley	Microsoft	USA	Projects:	Projects:
Pro	jects	5			DB e4	DB e4
	Ρ1	DB	Redmond	e4	Microsoft	Web e5
	P2	Web	S. Valley	e5	Projects:	
Em	ploye	es	5		Web e5	
	e4	John	×234		Employees	Employees
	e5	Anna	x888		e4 John	e4 John
					e5 Anna	e5 Anna



- **Goal:** infer a grouping function that has the same effect as the one intended by the designer
 - Muse-G probes each possible grouping attribute: start with cbranch

	Exa	mple	source		Target Scenario 1 group by cbranch	Target Scenario 2 do not group by cbranch
Cor	npan	ies			Orgs	Orgs
	Redr	mond	Microsoft	USA	Microsoft	Microsoft
	S. V	alley	Microsoft	USA	Projects:	Projects:
Pro	jects				DB e4	DB e4
	P1	DB	Redmond	e4	Microsoft	Web e5
	P2	Web	S. Valley	e5	Projects:	
Em	ploye	es	-		Web e5	
	e4	John	×234		Employees	Employees
	e5	Anna	x888		e4 John e5 Anna	e4 John e5 Anna



- **Goal:** infer a grouping function that has the same effect as the one intended by the designer
 - Muse-G probes each possible grouping attribute: start with cbranch

Fxa	mnle	source		Target Scenario 1	Target Scenario 2
		Jource		by cbranch	by cbranch
Compar	nies			Orgs	Orgs
Red	mond	Microsoft	USA	Microsoft	Microsoft
S. \	/alley	Microsoft	USA	<pre>Projects: SK(Redmond,y)</pre>	Projects: SK(y)
Projects	5			DB e4	DB e4
P1	DB	Redmond	e4	Microsoft	Web e5
P2	Web	S. Valley	e5	Projects: SK(S. Valley,y)	
Employe	ees	J		Web e5	
e4	John	×234		Employees	Employees
e5	Anna	×888		e4 John	e4 John
				e5 Anna	e5 Anna
				$v \subset \int Microsoft$	
				$y \subseteq 1$ microsoft, t	
					R
					() 9/

■ The next probed attribute is cname

Example source

Companies

S. V.	alley	Microsoft	USA
Mt.	View	Google	USA
Projects			
P1	DB	S. Valley	e4

P4	Web	Mt.	View	eб

Employees

e4	John	x234
еб	Kat	x331

■ The next probed attribute is cname

Examp	le source		Target Scenario 1 group by cname	Target Scenario 2 do not group by cname
Companies			Orgs	Orgs
S. Valley Mt. Viev	MicrosoftMicrosoft	USA USA	Microsoft Projects:	Microsoft Projects:
Projects	C		DB e4	DB e4
P1 DE P4 We	8 S. Valley eb Mt. View	е4 еб	Google Projects:	Web e6 Google
Employees			VVeb eb	Projects:
e4 Joh	n x234		Employees	DB e4
e6 Kat	: ×331		e4 John	Web e6
			e6 Kat	Employees e4 John
				e6 Kat

The next probed attribute is cname



Example source

^	•
Comp	anies
•	

S. ∖	/alley	Microsoft	USA
Mt.	View	Google	USA
Projects	5	-	
P1	DB	S. Valley	e4
P4	Web	Mt. View	еб
Employees			
e4	John	×234	
еб	Kat	x331	

Target Scenario 1 group by cname

Orgs	
Microsoft	
Projects:	
DB	e4
Google	
Projects:	
Web	e6

Employees e4 John еб Kat

Target	t Scer	nario	2
do	not gr	oup	
Dy	/ cnan	10	
Orgs			
Micr	osoft		
Proj	ects:		
	DB	e4	
	Web	еб	
Goog	gle		
Proj	ects:		
	DB	e4	
	Web	еб	
Employees			
e4	John		
еб	Kat		

The next probed attribute is cname

Example source

Cor	npan	ies		
	S. Valley		Microsoft	USA
	Mt.	View	Google	USA
Pro	jects			
	Ρ1	DB	S. Valley	e4
	P4	Web	Mt. View	еб
Employees				
	e4	John	×234	
	e6	Kat	x331	

Target Scenario 1 group by cname Orgs Microsoft **Projects**: DB e4 Google **Projects**: Web e6 **Employees** e4 John Kat e6

Target Scenario 2 do not group by **cname** Orgs Microsoft **Projects**: DB e4 Web e6 Google **Projects**: DB e4 Web e6 **Employees** John e4 Kat e6

e4

e6

The next probed attribute is cname

Example source

Companies USA S. Valley Microsoft Google USA Mt. View **Projects** S. Valley P1 DB Web Mt. View P4 **Employees** John x234 e4

x331 Kat e6

Target Scenario 1 group by **cname** Orgs Microsoft **Projects:** SK(Microsoft,y) DB e4 Google **Projects:** SK(Google,y) Web e6 **Employees** John e4 Kat e6

Target Scenario 2 do not group by **cnăme** Orgs Microsoft **Projects:** SK(y) DB e4 Web e6 Google **Projects:** SK(y) DB e4 Web e6 **Employees** e4 John Kat e6

 $y \subseteq \{ USA \}$

The next probed attribute is cname

Example source

Cor	npan	ies		
	S. Valley		Microsoft	USA
	Mt.	View	Google	USA
Pro	jects	5		
	P1	DB	S. Valley	e4
	P4	Web	Mt. View	еб
Employees				
	e4	John	×234	
	еб	Kat	×331	

Target Scenario 1	Target Scenario 2	
by cname	by cname	
Orgs	Orgs	
Microsoft	Microsoft	
Projects: SK(Microsoft,y)	Projects: <mark>SK(</mark> y)	
DB e4	DB e4	
Google	Web e6	
Projects: SK(Google,y)	Google	
Web e6	Projects: SK(y)	
Employees	DB e4	
e4 John	Web e6	
e6 Kat	Employees	
	e4 John	
	e6 Kat	
	•	

$y \subseteq \{ USA \}$

The wizard continues to probe the remaining possible grouping attributes

Running queries over the real source instance *I* **Example: probing on cname**

Running queries over the real source instance *I* Example: probing on cname

Query: Companies(c_1 , n_1 , l_1) \land

 $\mathsf{Projects}(\mathsf{p}_1,\mathsf{pn}_1,\mathsf{c}_1,\mathsf{e}_1) \ \land \\$

 $\mathsf{Employees}(\mathsf{e}_1,\mathsf{en}_1,\mathsf{cn}_1)$ \land

Running queries over the real source instance *I* **Example: probing on cname**

Query: Companies(c_1 , n_1 , l_1) \land Projects(p_1 , pn_1 , c_1 , e_1) \land Employees(e_1 , en_1 , cn_1) \land Companies(c_2 , n_2 , l_1) \land Projects(p_2 , pn_2 , c_2 , e_2) \land

 $\mathsf{Employees}(\mathsf{e}_2,\mathsf{en}_2,\mathsf{cn}_2)$ \land

Running queries over the real source instance *I* Example: probing on cname

Query: Companies(c_1 , n_1 , l_1) \land

 $\mathsf{Projects}(\mathsf{p}_1,\mathsf{pn}_1,\mathsf{c}_1,\mathsf{e}_1)$ \land

 $\mathsf{Employees}(\mathsf{e}_1,\mathsf{en}_1,\mathsf{cn}_1)$ \land

Companies(c_2 , n_2 , l_1) \land

 $\mathsf{Projects}(\mathsf{p}_2,\mathsf{pn}_2,\mathsf{c}_2,\mathsf{e}_2)$ \land

 $\mathsf{Employees}(\mathsf{e}_2,\mathsf{en}_2,\mathsf{cn}_2)$ \land

 $\mathsf{n}_1\neq\mathsf{n}_2$





Considering functional dependencies in the source can reduce the number of questions posed to the designer

- Considering functional dependencies in the source can reduce the number of questions posed to the designer
- Two mappings M_1 , M_2 have the same effect if for any source instance I, the result of exchanging I with M_1 is the "same" as the result of exchanging I with M_2

Homomorphically equivalent

- Considering functional dependencies in the source can reduce the number of questions posed to the designer
- Two mappings M_1 , M_2 have the same effect if for any source instance I, the result of exchanging I with M_1 is the "same" as the result of exchanging I with M_2

Homomorphically equivalent

Proposition. If a FD $A \to B$ holds, then a mapping M that groups by A has the same effect as a mapping M that groups by $A \cup C$, where $C \subseteq B$.

- Considering functional dependencies in the source can reduce the number of questions posed to the designer
- Two mappings M_1 , M_2 have the same effect if for any source instance I, the result of exchanging I with M_1 is the "same" as the result of exchanging I with M_2

Homomorphically equivalent

Proposition. If a FD $A \to B$ holds, then a mapping M that groups by A has the same effect as a mapping M that groups by $A \cup C$, where $C \subseteq B$.

- Suppose **cbranch** is a key, then we may save some questions
 - If the designer chooses Scenario 1 (including cbranch in the grouping function), probing on cname or location is no longer necessary

Proposition (Completeness). If there are n possible grouping attributes for a nested set S, then the questions asked by Muse-G explore the entire space of 2^n grouping functions. Muse-G asks at most n questions to infer the desired grouping semantics for S.

Proposition (Completeness). If there are n possible grouping attributes for a nested set S, then the questions asked by Muse-G explore the entire space of 2^n grouping functions. Muse-G asks at most n questions to infer the desired grouping semantics for S.

Proposition (Small examples). At each probe, Muse-G constructs a source example of size at most twice the number of conjuncts in the for clause of the mapping.

Proposition (Completeness). If there are n possible grouping attributes for a nested set S, then the questions asked by Muse-G explore the entire space of 2^n grouping functions. Muse-G asks at most n questions to infer the desired grouping semantics for S.

Proposition (Small examples). At each probe, Muse-G constructs a source example of size at most twice the number of conjuncts in the for clause of the mapping.

Incremental design: group more or less starting from an existing grouping function

Proposition (Completeness). If there are n possible grouping attributes for a nested set S, then the questions asked by Muse-G explore the entire space of 2^n grouping functions. Muse-G asks at most n questions to infer the desired grouping semantics for S.

Proposition (Small examples). At each probe, Muse-G constructs a source example of size at most twice the number of conjuncts in the for clause of the mapping.

- Incremental design: group more or less starting from an existing grouping function
- Design for a specific source instance: reduce the number of questions
 - Muse-G identifies attributes whose inclusion/exclusion as arguments of grouping functions is inconsequential

Ambiguous Mappings

CompDB: Rcd **Projects:** Set of **Project:** *Rcd* pid pname manager tech-lead Employees: Set of **Employee:** *Rcd* 🗲 eid ename contact

Ambiguous Mappings

CompDB: Rcd **Projects:** Set of **Project:** *Rcd* pid pname manager tech-lead **Employees:** Set of **Employee:** *Rcd* 🗲 eid ename contact

OrgDB: Rcd Projects: Set of Project: Rcd pname supervisor email




for

p in CompDB.Projects e₁ in CompDB.Employees e₂ in CompDB.Employees

satisfy

 $e_1.eid = p.manager$ $e_2.eid = p.tech-lead$



for

p in CompDB.Projects e₁ in CompDB.Employees e₂ in CompDB.Employees

satisfy

 $e_1.eid = p.manager$ $e_2.eid = p.tech-lead$

exists

 p_1 in OrgDB.Projects



for

p in CompDB.Projects e₁ in CompDB.Employees e₂ in CompDB.Employees

satisfy

 $e_1.eid = p.manager$ $e_2.eid = p.tech-lead$

exists

p₁ in OrgDB.Projects

where

 $p.pname = p_1.pname$



for

p in CompDB.Projects e₁ in CompDB.Employees e₂ in CompDB.Employees

satisfy

 $e_1.eid = p.manager$ $e_2.eid = p.tech-lead$

exists

 p_1 in OrgDB.Projects

where

 $p.pname = p_1.pname$

 p_1 .supervisor =

 $e_1.ename$ or $e_2.ename$



e₁.contact **or** e₂.contact

- This mapping is ambiguous
- There are four alternative interpretations

e1.enamee1.enamee2.enamee2.enamee1.contacte2.contacte1.contacte2.contact

Key idea: provide an example that illustrates the alternative interpretations in a compact way

Projects

P1 DB e4 e5 Employees e4 John john@ibm e5 Anna anna@ibm

Projects	Orgs		
P1 DB e4 e5	Project	S:	
Employees	DB	John	iohn@ibm
e4 John john@ibm		50111	Jerneism
e5 Anna anna@ibm		Anna	anna@ibm

Key idea: provide an example that illustrates the alternative interpretations in a compact way



The mapping designer makes one choice for each ambiguous element
 Each decision removes one ambiguity



- The mapping designer makes one choice for each ambiguous element
 Each decision removes one ambiguity
 - E.g., choosing "Anna" as the supervisor and "john@ibm" as the email



- The mapping designer makes one choice for each ambiguous element
 Each decision removes one ambiguity
 - E.g., choosing "Anna" as the supervisor and "john@ibm" as the email

```
\begin{array}{c|c} p_1. \text{supervisor} = & & p_1. \text{email} = \\ e_2. \text{ename} & e_1. \text{contact} \ \textbf{or} \ e_2. \text{contact} \end{array}
```



- The mapping designer makes one choice for each ambiguous element
 Each decision removes one ambiguity
 - E.g., choosing "Anna" as the supervisor and "john@ibm" as the email

```
\begin{array}{c|c} p_1. \mathsf{supervisor} = & & p_1. \mathsf{email} = \\ e_2. \mathsf{ename} & & e_1. \mathsf{contact} \end{array}
```

Running queries over the real source instance

Running queries over the real source instance

Query:

Projects(p_1 , pn_1 , e_1 , e_2) \land Employees(e_1 , en_1 , cn_1) \land Employees(e_2 , en_2 , cn_2) \land

Running queries over the real source instance

Query:

 $\begin{aligned} \mathsf{Projects}(\mathsf{p}_1,\mathsf{pn}_1,\mathsf{e}_1,\mathsf{e}_2) \land \\ & \mathsf{Employees}(\mathsf{e}_1,\mathsf{en}_1,\mathsf{cn}_1) \land \\ & \mathsf{Employees}(\mathsf{e}_2,\mathsf{en}_2,\mathsf{cn}_2) \land \\ & \mathsf{en}_1 \neq \mathsf{en}_2 \land \mathsf{cn}_1 \neq \mathsf{cn}_2 \end{aligned}$

Running queries over the real source instance



Running queries over the real source instance



Muse-D: Properties

 For each ambiguous mapping, the designer is presented with a single example

Muse-D: Properties

For each ambiguous mapping, the designer is presented with a single example

Proposition (Completeness). The single example differentiates among all the alternative interpretations of the ambiguous mapping. The mapping designer has to make a number of choices equal to the number of ambiguous elements.

Muse-D: Properties

For each ambiguous mapping, the designer is presented with a single example

Proposition (Completeness). The single example differentiates among all the alternative interpretations of the ambiguous mapping. The mapping designer has to make a number of choices equal to the number of ambiguous elements.

Proposition (Small examples). The number of tuples in the example source instance is the number of conjuncts in the for clause of the mapping.

Experiments: Setting

Mapping	Size of	Sets with	Number	Ambiguous	Alternative
Scenarios	real source	refinable	of	mappings	interpretations
	instance	grouping	mappings		
Mondial	1MB	8	26	7	208
DBLP	2.6MB	6	4	0	4
ТРСН	10MB	4	5	1	20
Amalgam	2MB	2	14	0	14

- Mapping system: Clio
- Query evaluation: DB2 v9, Saxon-B 8.9
- MUSE implementation: Java 6

Experiments: Muse-G

Mapping	Average #	Number of	% times	Average time	Grouping
scenario	of grouping	questions	found real	to get real	strategy
	attributes	(average)	example	example (s)	
		2.6	38%	0.014	G_1
Mondial	13.1	8.5	41%	0.187	G_2
		2.9	40%	0.015	G_3
		1.5	17%	0.450	$\mathbf{G_1}$
DBLP	11	11	11%	0.337	$\mathbf{G_2}$
		1.5	17%	0.454	$\mathbf{G_3}$
		1.5	0%	0.785	G_1
ТРСН	26.7	17	12%	0.893	G_2
		1.5	0%	0.782	G_3
		2	29%	0.013	G_1
Amalgam	14.1	3	52%	0.043	G_2
		3	52%	0.030	G_3

 G_1 : group by all possible attributes (cbranch, cname, location, pid, ...)

- G_2 : group by all attributes exported above the set (cname)
- G_3 : group by all exported attributes (cname, pname, eid, ename)

Experiments: Muse-D

Mapping	Alternatives	Number of	Size of	Number of
Scenario	encoded	questions	source example	ambiguous values
			(# of tuples)	in target instance
Mondial	208	7	3–4	4–5
ТРСН	16	1	9	4

Related Work

There are many related works, we focus here on the closest one
 We were inspired by DataViewer [Yan et al. 01]

Feature \setminus System	MUSE	DataViewer	
Examples	Compact	As many as	
to be analyzed	representation	alternative interpretations	
	(e.g. 7)	(e.g. 208)	
Grouping Semantics	Yes	No	
Data Models	Relational and XML	Relational	
Mapping	Source-target	SQL	
Language	mappings (GLAV)		

Conclusion

- MUSE: a mapping design wizard
 - Use examples to understand, design, refine schema mappings
 - Work with complete and small data examples rather than complex specifications
 - Focus on two important aspects of a mapping specification
 - Grouping semantics
 - Interpretation of ambiguous mappings

Thank you